

LANGUAGE DESCRIPTION for FRONTEND IMPLEMENTATION

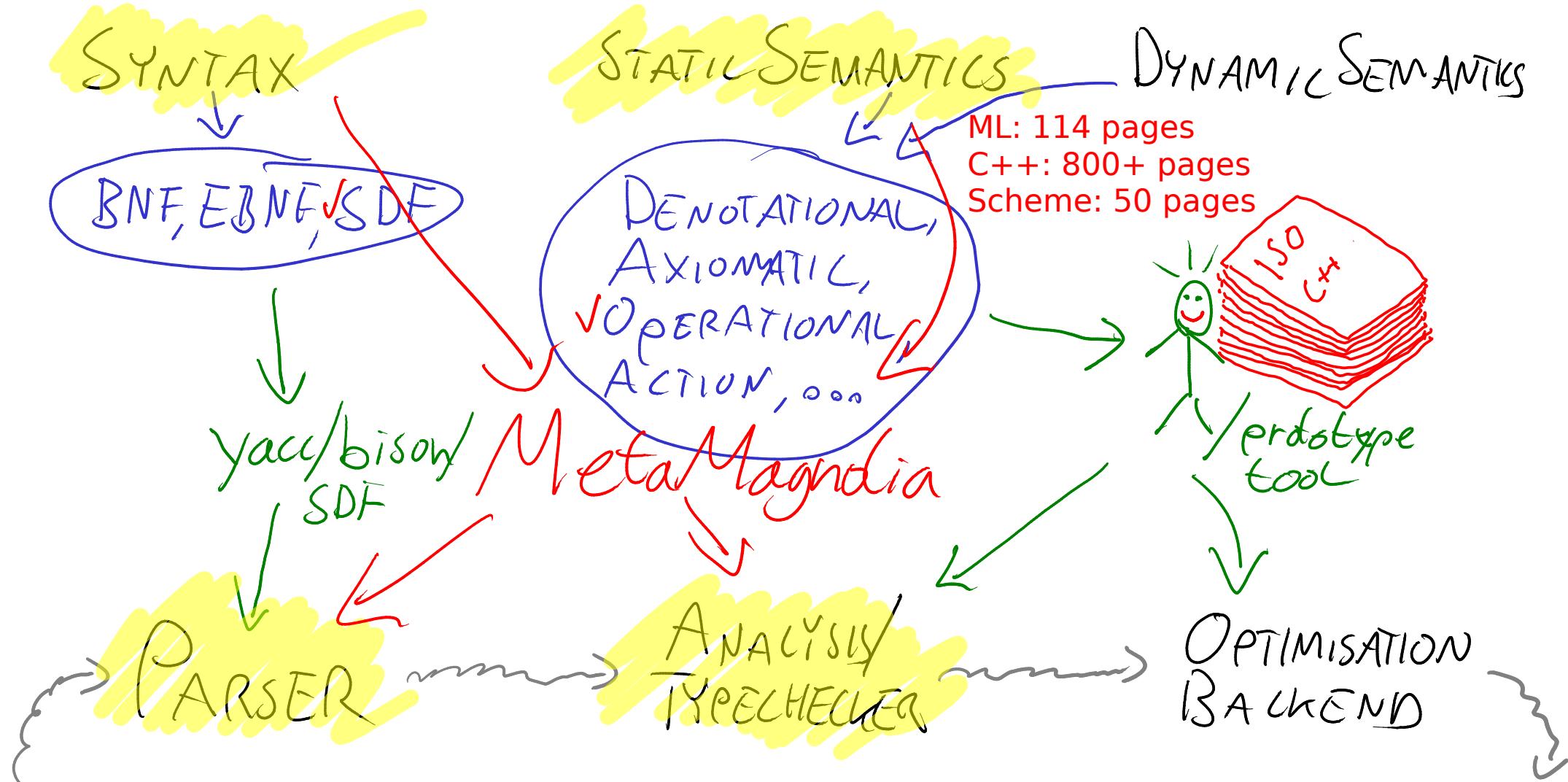
ANYA HELENE BAAGA

BERGEN LANGUAGE DESIGN LABORATORY
DEPT. OF INFORMATICS
UNIVERSITY of BERGEN

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LANGUAGE SPECIFICATION



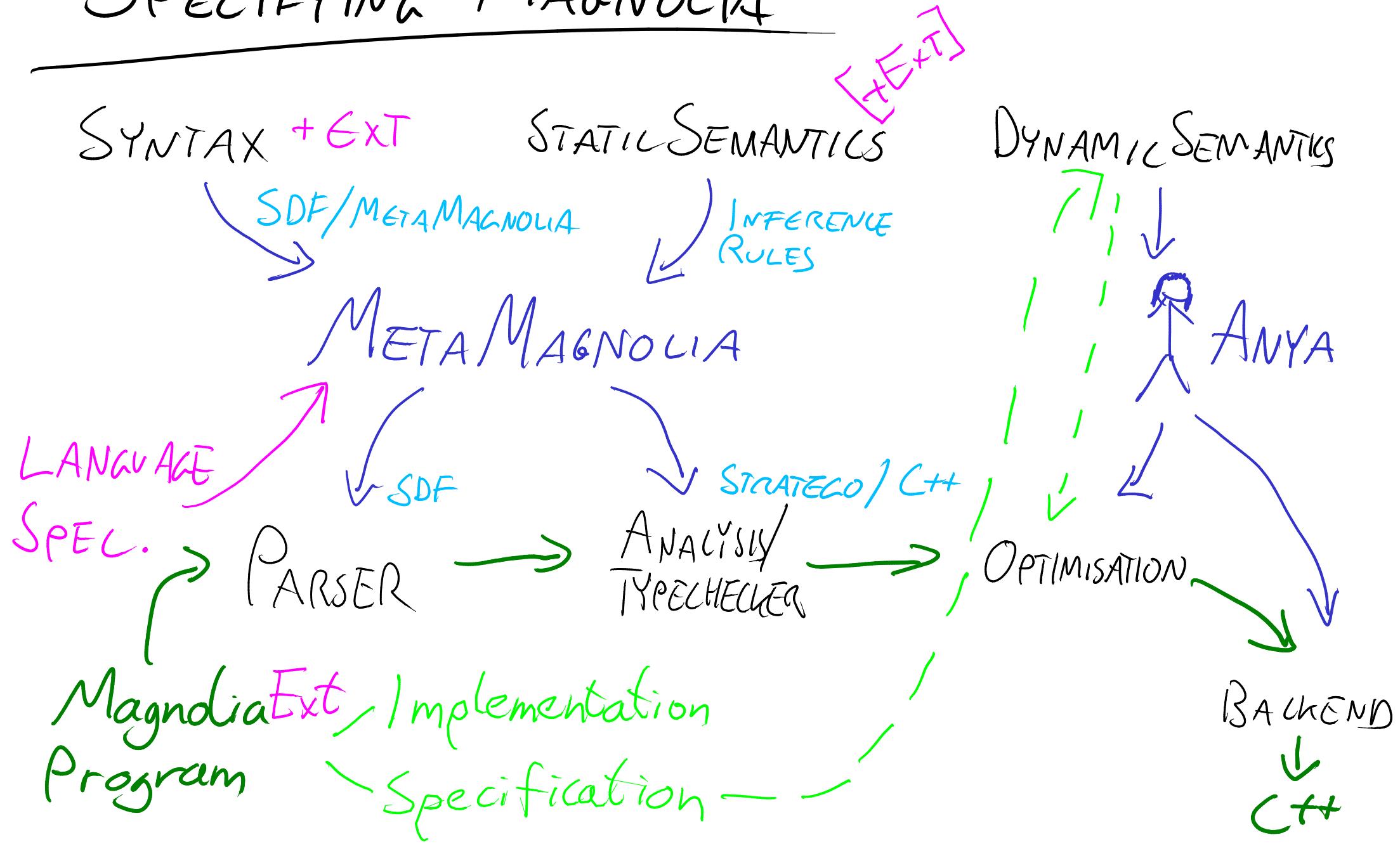
THE MAGNOLIA LANGUAGE

- STATIC TYPING, NO INFERENCE
- OVERLOADING — name/overload resolution
- IMPLICIT DECLARATIONS [LDTAO9]
- ALGEBRAIC SPECIFICATION
- SIMILAR TO C++
(well, somewhat...) *concepts/
typed classes
+
axioms*

EXPERIMENTAL!

- EVOLVING DESIGN
- LANGUAGE VARIANTS
- EXTENSIONS

SPECIFYING MAGNOLIA



METAMAGNOLIA SPECIFICATIONS

SYNTAX

CORE

STATIC
SEMANTICS

ANNOTATED TREE

- DEFINED IN SDF
OR METAMAGNOLIA

syntax mix Magnolia/Core
+ Magnolia/BaseExt

- SUGAR RULES
TO/FROM CORE

if: +«if e then s1* end»
 <-> <if e then s1* else end>

- PRETTY-PRINT BY
EXAMPLE

«if e then
 s1*
else
 s2*
end»

- PATTERN ABSTRACTION

patterns
cond(c, t, e)

- INFERENCE RULES

$e \implies e'$ $\text{isPredicate}(e')$
 $s1^* \implies s1'^*$ $s2^* \implies s2'^*$

«if e then s1* else s2* end»
==> <if e' then s1'^* else s2'^* end>



SUGAR EXAMPLE

language module Magnolia

syntax mix Magnolia/Core + Magnolia/BaseExt

syntax

+« n » -> < n:? >

« e1 o e2 » <-> < _o_(«e1», «e2») >

« o e1 » <-> < o_(«e1») >

« e[es*] » <-> < index(«e», «es*») >

+« if e then s1* end » <-> < if «e» then «s1*» else end >

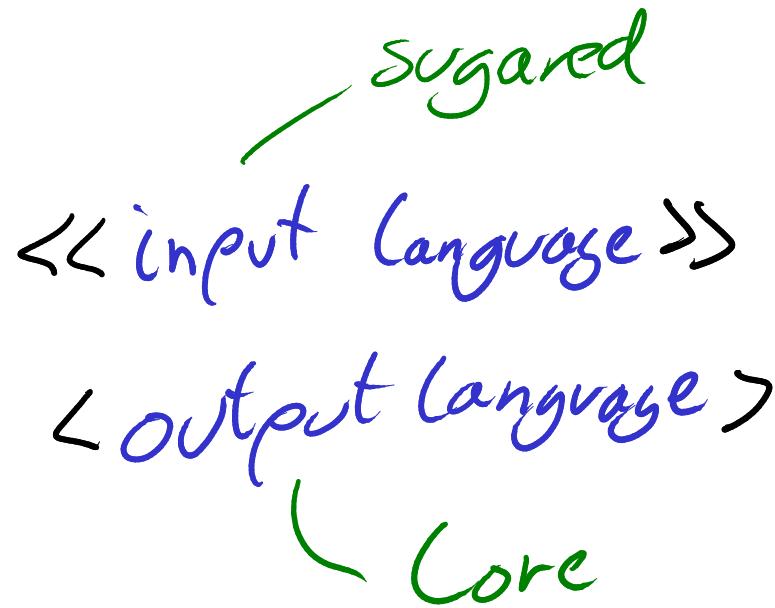
// C++-like syntax:

+« if(e) s1 » {prefer} -> < if «e» then «s1» else end > // amb

+« if(e) s1 else s2 » -> < if «e» then «s1» else «s2» end >

+« while(e) s » -> < while «e» do «s» end >

Quotin^g



$\frac{+}{+}$

New
Syntax

Generated
in SDF

WHAT HAPPENS?

- TRANSFORMATION CODE (Strategy / C++)
- BANANA SAFETY
 - COMPLETE TRANSFORMATION
 - TERMINATION
 - EFFICIENCY

implicit
recursion

[Anderson, Bradford '89]



DEFINING STATIC SEMANTICS

SOS / MSOS / I-MSOS

~~Attribute grammars~~

~~Transformation rules~~

$$\frac{\cancel{f : T \rightarrow f'} \quad \cancel{e : T \rightarrow e'}}{\cancel{f(e)} \Rightarrow f'(e')}$$

ABSTRACTION

MODULARITY [MSOS] — pattern abstraction
constructs

IMPLICIT [I-MSOS] — handling inessential features
(SIMILAR TO MSDF)



Example

construct let statement:

syntax

«let var n : t = e; in s* end»

static semantics

$e \implies e' : t'$ $s^* \implies \text{declare}(<\text{var } n : t' = _ ;>) \implies s'^*$

«let var n : ? = e; in s* end»
 $\implies <\text{let var } n : t' = e' ; \text{ in } s'^* \text{ end}>$

OPERATIONS

USE

$$p \Rightarrow p'$$

$$p - op \rightarrow p'$$

$$p - op(\dots) \rightarrow p'$$

$$op(p, \dots) = p'$$

DEFINITION

$$\frac{p_1 \dots p_n}{p \Rightarrow p'}$$

$$\frac{p_1 \dots p_n}{p - op(\dots) \rightarrow p'}$$

OR: EXTERNAL DEF
(WITH AXIOMS)

EXAMPLE

construct if statement:

static semantics

isPredicate(«e»)

«if e then s1* else s2* end»
==> <if «e» then «s1*» else «s2*» end>

construct while statement:

static semantics

«while e do s* end» ==> <while «e» do «s*» end>

OVERLOAD RESOLUTION

construct call statement:

syntax

«call p(as*);»

static semantics

as* ==> as'*:ts'*
p --?-callable(ts')*-best-unique-> p'

«call p(as*);» ==> <call p'(as'*);>

CALLABILITY

operator callable?:

 compatible(t1*, t2*) = x

procedure p(t1*); --callable?(t2*)-> p weight x

operator compatible?:

 compatible?(t, t) = 0

 not(equal(t1, t2)) nameOf(t1)=n n(?:t1) ==> (_, t2)

 compatible?(t1, t2) = -10

operator compatible = reduce(compatible?, 0)

HANDLING ERRORS

NEED ALL INFORMATION

Avoid CUTTER!

IMPLICIT Prop of CONTEXT INFO

ATTACH ERROR TO "SUCCESSFUL" RESULT

DIFFERENCE BETWEEN

NOT APPLICABLE
SEMANTIC ERROR

Conclusion — METAMAGNOLIA

SYNTAX → COMPOSITION (with SDF)
→ EXTENSION
→ DESUGARING

SEMANTICS — STATIC
→ MODULES, ABSTRACTION, ADTs
→ ERROR HANDLING
→ IMPLICIT PROPAGATION/RECUSION
→ DEAL WITH OVERLOADING!

USE — MAGNOLIA

? — NEED NEW NAME...

Check out our new Language
Design Lab!

<http://blndl.ii.vib.no/>

we're looking for
PhD students!

Photos from last
night:

<http://www.ii.vib.no/~nanya/photos/Ldta-2010/>